$\mathrm{AntCrypt}^{\star}$

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^{*} The name "AntCrypt" was chosen because the core of our construction resembles an anthill: in both, a huge quantity of small workers carry our tiny tasks in apparent chaos, however, in reality this "chaos" is orchestrated so that all results come together and form the final result.

1 Introduction

Arguably the biggest threat to password hashing schemes stems from GPUs, FPGAs, and ASICs, who provide enormous computing power which can speed up verification of a batch of passwords (e.g., in an offline guessing attack). Common constructions for password hashes use, at their core, two different methods to limit speed-up of verification operations.

- First, aggressively iterated constructions proportionally increase the computation times for verification on all platforms. The (well-understood) problem with constructions solely relying on iterated constructions is that they are typically quite fast when implemented on GPUs and FPGAs, as they can be parallelized very well.
- Frequent memory access (e.g., memory-hardness and similar ideas) are intended to slow down implementations on hardware basically utilizing memory bandwidth and memory latency. Large memory requirements (such as scrypt) will force the attacker to access main memory (on GPUs), while moderate memory usage (such as bcrypt) leaves the attacker with a trade-off between using a large number of registers and thus voiding memory access, or using fewer registers but accessing global memory.

One concern with berypt is that the size of the memory used in the computation is fixed to 4 kByte and cannot be changed, and that 4kByte is potentially not enough memory to effectively thwart efficient implementations on FPGAs. With scrypt, one concern is that the huge memory requirements are problematic if deployed on servers handling frequent login requests, and make the server susceptible to denial-of-service attacks. Another potential concern is that memory access is "relatively rare" in the sense that there is one hash function computation between two memory access operations.

In our proposal, we opt for a middle-ground, which seems to offer the best benefit of both worlds:

1. Memory usage: We use moderate amounts of memory, tunable with a parameter from 256 Bytes upwards, where a reasonable choice seems to be around 32 kBytes. We ensure very frequent access to all regions of the memory, similar to berypt and different from scrypt, to avoid previously mentioned potential problems.

In addition, our construction makes use of a (to the best of our knowledge novel) idea that aims to slow down implementations on GPUs and FP-GAs/ASICs specifically.

2. Control-flow divergence: Our code will frequently branch depending on the current state (and thus ultimately on the password), to (i) avoid good parallelization on GPUs, and (ii) increase the size of implementations (and thus increase the cost and decrease the throughput) on FPGAs/ASICs.

2 The Key-Derivation Scheme

In this section, we will provide a description of our construction and comment on the design choices. A more detailed discussion will follow in Section 3.

2.1 Parameters and Main Data Structure

Unless stated otherwise, all data types are 32-bit words. As the prototype of the PHS function provides two cost parameters, we derive the internal parameters from them as follows:

 state_bytes defines the amount of memory used for the state in bytes and is defined as

$$state_bytes = 2^{m_cost+8}$$
.

Analogously, state_words defines the number of 32-bit words the state contains.

 inner_rounds defines the number of iterations for the inner loop, iterating over all state_words state positions. We require a minimum of at least two inner rounds as follows:

$$\mathtt{inner_rounds} = \max(\left\lfloor \frac{\mathtt{m_cost}}{16} \right\rfloor, 2)$$

outer_rounds defines the number of iterations for the outer loop. We require
at least one outer round and define it as follows:

$$outer_rounds = max(t_cost, 1)$$

The *primary data structure* is a memory buffer buf = [prefix, memory]. The prefix can be used as a to generate different hash values from the buffer. If not stated otherwise, we refer to buf as the memory without the prefix.

In the algorithm, we use two such buffers: state of size state_words+1 32-bit words, as well as a rehash buffer of 16+1 words. The size of rehash is equivalent the output length of the *primary hash function* + 1 word as a prefix. These buffers are accessed either on byte level as bytes $[0,1,\ldots,(\text{state_bytes}+3)]$ or as words $[\{3,2,1,0\},\{7,6,5,4\},\ldots,\{\ldots,(\text{state_bytes}+3)\}]$.

2.2 Algorithm Definition

Algorithm 1 describes the basic structure of the derivation function. In the following, we will describe the main functions init, update_entropy, update_state and compute_output in more detail.

Algorithm 1 Pseudocode of AntCrypt

```
Require: t_{cost} > 0, m_{cost} > 0, outlen > 0, salt, pw,
Ensure: key
1: init(salt, pw)
                                                                   {Initialize state}
 2: for i = 0 to outer_rounds do
      update_entropy()
                                                  {Distribute entropy over the state}
      # The following loop is referred to as update_state()
 5:
      for j = 0 to inner_rounds do
                                                   {Waste time operating on state}
 6:
        int_update_state()
 7:
      end for
8: end for
                                                       {Final output transformation}
9: compute_output()
```

Description of init(): The initialization function init() fills the empty state memory with its initial content and is implemented in the reference implementation as the function phs_init(). Please note this function addresses the memory byte-by-byte, not as 32-bit words.

- 1. The salt is copied to the beginning of the state memory. It is interpreted byte-wise, i.e., state[0] = salt[0], state[1] = salt[1], ...

 We use a fixed size for the salt (16 bytes as suggested in the proposal), and do not see a reason for supporting variable sized salts: 16 bytes (128 bit) should offer sufficient security against appropriate attacks and we do not need to add any separator or length into the buffer when using a fixed length.
- 2. The password is appended to the array after the salt, also stored byte-perbyte. The maximum length of a password accepted is $\mathtt{state_bytes} 16 1$ bytes to leave enough space for the salt, the password and an end-identifier. The minimal state supported consists of 256 bytes. Thus, a 128 byte password as required will always be accepted.
- 3. A "password-end" marker 0x80 is appended directly after the password and the remaining space is filled with 0x00.

Description of update_entropy() The function update_entropy (which maps to phs_upd_entropy() in the reference implementation) uses a hash function, hashing the entire state. As common hash functions have a much smaller output compared to state – e.g., 128 bit for MD5 or 512 bit for SHA-512 – we need to extend these constructions to adapt for the larger output size. In the implementation, we use the rehash buffer and its prefix to derive the new state.

Similar constructions are well-known in the cryptographic literature, and in the random oracle model it is easy to prove that the resulting function constitutes a secure hash function. Basically, we compute

$$h := H(\text{state}),$$

and then

$$s_i = H(i \parallel h)$$

```
s_0, s_1, \ldots, s_k.
```

The "intermediate" hash value h is also used in the compute_output() function for an additional feature.

Description of update_state() The function update_state() accesses the buffer state inner_rounds × state_words times and aims at wasting CPU cycles and efficiently slow down parallel computation on different platforms. In the reference implementation, this function is implemented as the function phs_upd_state().

Algorithm 2 Pseudocode of update_state()

```
1: for i = 0 to inner_rounds do
 2:
      for j = 0 to state_words do
 3:
         res = (state[j] ROR i)
 4:
         tgt\_addr = res \% state\_words
         reset idx permutation
 5:
         for j = 0 to #F do
 6:
 7:
            choose unused idx by evaluating res
 8:
            \mathrm{res} = F_{\mathrm{idx}}(\mathrm{res})
 9:
         end for
         state[tgt_addr] = state[tgt_addr] XOR res
10:
       end for
11:
12: end for
```

Algorithm 2 describes the update algorithm. We use a set of functions $F_i(x)$, where #F is the number of functions and use a calling sequence of these functions, where every function is called exactly once. After all #F functions process the data, the word at the target address is updated by using a bit-wise XOR.

Please note that the sequence is not pre-defined, but depends on the value res (and thus the initial value state[j]). Thus, in theory, all #F! sequences are possible.

The currently implemented functions (defined in phc.h) are given in Table 1. Please note that the functions are currently being evaluated and may be tweaked later.

Description of compute_output() The function compute_output() uses the state memory after the last outer round to generate the derived key material. It is implemented as phs_gen_output() in the reference implementation.

It consists of two steps, depending on the requested output length. If the output length is less or equal to 512 bit, i.e., the output length of SHA-512, only the first step is necessary.

```
/* integer operations */
#define F00(X) ( (X) + 0x01234567 )
#define F01(X) ( (X) * 0x89ABCDEF )
/* bit operations */
#define F02(X) ((X) >> 3)
#define FO3(X) ( ROTR((X), 7) )
#define F04(X) ( (X) ^{\circ} 0x01234567 )
#define F05(X) ( (X) & OxFEFEFEFE )
#define F06(X) ( (X) | 0x02020202)
/* floating point operations */
#define F07(X) ( (uint32_t) ( 2147483648.L \
                * sin (((double) X)/1000000000.L )) )
#define F08(X) ( (uint32_t) ( 2147483648.L \
                * cos (((double) X)/1000000000.L )) )
#define F09(X) ( (uint32_t) ( 2147483648.L \
                * tan (((double) X)/5000000000.L )) )
/* 1/x: [1,2] \rightarrow [0.5, 1] (bijective) */
#define F10(X) ( (uint32_t) ( (double) ( 2 * 4294967296.L \
                * ( 1 / (1.5 + (double) X / 4294967296.L )) - 0.75 ) ) )
```

Table 1. List of the functions F_i used in update_state().

First, we generate the intermediate hash, which would be generated during the next call to update_entropy. It basically is identical to the first step of update_entropy(), i.e., we apply the hash function to the entire state:

$$h := H(\text{state}).$$

Depending on the desired output length, we use up to 64 byte from h, addressing the buffer byte-wise and starting with byte 0.

In case more than 64 byte were requested, we use the prefix for the state, initialized with 1, and hash the full state including the prefix to derive a new intermediate value

$$h' := H(i \parallel state).$$

We use the same function used in update_entropy() to derive a new state from h', overwriting the previous state and append up to state_bytes bytes to the output. This procedure can be repeated up to $2^{32} - 1$ times, effectively producing more than $2^{40+\text{m}_\text{cost}}$ bytes of key material.

This construction has another advantage: By storing the "intermediate" value h as final output, we are able to recompute the "next" state. This means that we can "resume" the computation of the state from a previously stored hash value, i.e., we can retroactively increase the hardness with respect to an increased parameter t_cost (cf. Section 3 for more details).

3 Design choices and remarks

Next, we comment on some of the design choices that underly our construction.

3.1 Implementation

One of our main intentions was to keep the overall structure and design as simple as possible, as this facilitates analysis and implementations. This also means we omitted some features from the implementation that are easy to add for a future (reworked) version. For the same reason we omitted most optimizations of the implementation and provide a rather straightforward implementation which is presumably easy to analyze. If selected for the second round we would provide an optimized version. The overall structure is very simple, with a clear distinction between the "cryptographically hard" step (update_entropy()), where we use established cryptographic primitives, and the "computationally hard" step (update_state()), where we are relatively free to do arbitrary computations that achieve our goals.

Some features that can easily be added (and will be added in future versions):

Parallelism There is a very easy modification to make the computation parallelizable for the honest server that computes the hash. Instead of processing each cell individually when computing the update_state(), we can read several (for example 16) consecutive cells, compute their output in parallel, and then write back simultaneously. (As we XOR the result on the target cell the order of writing does not matter.) This provides sufficient parallelism for the honest server, while not being advantageous for the attacker, as these parallel threads are still diverging.

Extending hardness Without further modifications, the above construction allows the legitimate server to increase the hardness of an existing hash without knowledge of the password, within certain constraints. It is necessary that the intermediate hash is stored in it entirety, i.e., the output has at least 512 bit. Furthermore, only the t_cost parameter can be increased (i.e., internally the outer_rounds parameter), the m_cost parameter needs to be fixed. Increasing the strength is very straightforward (and we will make code for doing so available in the near future). As the final step compute_output() is equivalent to the first part of the update_entropy() step (for an output length of 512 bit), we can simply resume the computation from this step on by first completing the second half, i.e., populating the entire state buffer from this value and then resuming with the iterations, adding so many iterations that the wanted iteration count is met, and finalizing with the final hashing.

3.2 Divergence and choice of the functions F_x

The specific choice of the functions F_i used in the construction depends on a number of factors, including the attacker's compute architecture. We are still

evaluating different choices for these functions, so the currently selected functions are likely to change in future versions; any comments are appreciated.

Some important considerations are the following: If the functions take too long to compute, then an attacker can potentially queue them up to compute the same ones in parallel, thwarting the divergence of the threats. However, if they are too fast to evaluate, then the "overhead" imposed by the computations in the inner loop that are not part of the F_i 's, e.g., computing the permutation, reduces the effectiveness of the divergence. (As computing the permutation incurs some substantial overhead, we consider using just a random sequence of indices; choosing a permutation, however, has the desirable property to rule out a number of timing side channel attacks as discussed later in this text.)

The overlap between different functions F_i , i.e., the potential to execute them in parallel, should be minimized; we attempted to achieve this by choosing functions with distinct assembler instructions, additionally ensuring that they are not easily transformable into each other. (Note the absence of the "bitwise invert" function, which can be expressed with an XOR.)

On using floating point operations We are aware that using floating point operations in such constructions is unusual, but we believe that they are helpful in minimizing the overlap, and they are also quite costly to implement on FP-GAs and ASICs. We avoid rounding errors by converting back each result to an integer, thus being able to control any potential rounding error. But again, the specific choice of the F_x is still somewhat experimental, and we might opt to remove floating point instructions if they incur problems with portability.

4 Security

4.1 Cryptographic security

Our construction inherits its cryptographic strength quite directly from the security of the underlying hash function. We describe our construction using SHA-512, but it can be easily substituted with any other hash function with sufficiently large state/output size. We have selected SHA-512 as it is a widely accepted design which has proven security over several years, and implementations are easily available in common libraries. In fact, it should be straight-forward to prove (in the random oracle model) that, provided that the functions F_x are permutations, or at least behave "sufficiently random", then the overall construction behaves like a random function.

In general, constructing secure hash functions is a delicate matter, and large efforts have gone into the design of such functions. Therefore, we feel that it is mandatory to rely on well-established constructions to achieve cryptographic security instead of attempting to use home-made constructions. One of the beauties of our construction is that we separate the task of providing cryptographic strength from the task of slowing down verification, (cryptographic strength is largely realized by the re-hashing done in update_entropy, putting minimal

requirements on update_state only, while the slow-down is largely realized in update_state).

The only thing that is required to really inherit these properties is that the applying the step $update_state$ does not loose too much entropy. However, by our construction, applying the sequence of the F_x to the current state is always a permutation, as we XOR the output to the target value. (This is very similar to the well-known Feistel structure, which also always is a permutation for arbitrary round functions.) And if one application of the sequence is a permutation, than by repeating this argument, the entire function $update_state$ constitutes a permutation.

4.2 Speed up

The intended use of function update_state is to slow down the computation of the password hash, thus this is the critical place to look for optimizations.

On CPUs, we believe that only minor optimizations can be done. The pseudo-random nature of the order of applying the functions F_x means that there is very little (constant) structure that can be exploited for optimizations. Note that, when looking retroactively at one particular run, there will be structure that can potentially be exploited, however, as the structure changes for each application of the permuted chain of F_x 's such structure needs to be detected during runtime. As the functions F_x are very short (ranging from a single assembler instruction up to a few), we believe that code for detecting and exploiting such structure would most likely slow down the execution more than it helps in speeding up.

Also note that we plan to consider other functions F_x in the future, and we hope to be able to provide a more formal argument regarding the potential speed up in future versions of this document.

On GPUs On GPUs, these random permutations will lead to a substantial amount of branch divergence, which means that the parallel executions of the hash function for a parallel brute force attempt (running for different passwords) will have divergent control flow. For "ideal" functions F_x with no overlap, no overhead outside the F_x , and ideally random selection, we would expect a slow-down equal to the number of functions, i.e., by a factor of 16. (Here "slowdown" is comparing the runtime for the case with convergent execution, e.g., all threads hashing the same password, with the runtime for divergent threads, e.g., when hashing different passwords in each thread.) In practice, there is overhead, e.g., caused by the final XOR and the computation of the permutation, and the functions have not entirely disjoint assembler instructions (e.g., we need some re-scaling of the values for the floating point instructions), so these ideal goals will likely not be met.

On FPGAs/ASICs While FPGAs and ASICs are very dangerous in terms of efficient implementation of brute-force attacks, the construction was chosen to render dedicated hardware attacks almost useless.

While many of the functions are easily mapped to hardware, floating point operations come at a high price. We analyzed the available cores for Xilinx Spartan 6, Virtex 6 and Artix 7 devices**. The CORDIC-core offers sine and square-root with 8 to 48 bit operands. For 32-bit operation, the minimum area is 3664 LUTs and 3588 FFs for sine (Spartan-6) and 975 LUTs and 1202 FFs for square-root (Artix-7) and has a latency of more than 32 clock cycles.

The floating-point core provides addition/subtraction, division, square-root and multiplication with configurable latency (time-area tradeoff) and may use available DSP cores, and the area consumption is heavily dependent on these configurations.

The use of more than one FPU function will significantly increase the area and latency of the generation on FPGAs. In addition, to support all possible sequences of the F_i functions, the complexity of the routing will increase dramatically: Every output needs to be routed to every other function as input. Thus, every function needs a large multiplexer, increasing the routing delay and increasing the critical path.

The second limiting factor is the memory usage, as fast memory cores are available but limited in size and number. To implement a 64 kByte state (m_cost = 8) will already use about 29 18k-BRAMs on Xilinx FPGAs. Thus, the memory area will become a limiting factor even with medium state sizes.

In practice, we think that using FPGAs or producing dedicated ASICs will not be the first choice for an attacker, as the construction is by design very cumbersome to implement and artificially adds latency, enforces complex routing and needs area-consuming FPU arithmetic.

4.3 Side channel attacks

Storing the password in memory The password is written to the state in the beginning and immediately overwritten by the output of the hash function. No copy needs to be stored beyond the initialization of the state memory. This should effectively prevent reading the password from memory.

(Cache) timing attacks The different functions F_x have usually, depending on the platform, different execution times. This could lead to timing attacks or to cache timing attacks. However, as we always use the same functions, just in differently permuted order, the time between memory access is constant (assuming that each operation runs in time independent of the data). In other words, the execution time between memory access is constant, thus no information is leaked. Then also the overall running time is constant, and no timing leak exists.

Other side channel attacks More involved side channel attacks, such as power consumption and electromagnetic emanation, are outside the scope of our consideration, as they depend on the specific architecture the code is running on. Also, they do not seem appropriate for the case considered, as an adversary that

 $^{^{\}star\star}$ cf. Xilinx DS858 and Xilinx DS335 specification

has physical access to a machine verifying the correct password typically has easier attacks at hand.

5 Final remarks

5.1 Statements

We ensure that we have not inserted, and are not aware of, any deliberately hidden weaknesses in the scheme described above.

The scheme is and will remain available worldwide on a royalty free basis, and we are unaware of any patent or patent application that cover the use or implementation of the submitted algorithm.

6 Test vectors

In this section, we will provide several testvectors. Please note that the requested list of $2^8 \times 2^8$ password/seed combinations for meaningful cost factors will take a very long time to generate. To do this, please compile the included source code and run the program phc_tv.

The format of the output is similar to previously used formats:

\$<seed>\$<t_cost>:<m_cost>\$<hash>

where the cost parameters are represented as two-digit numbers and the seed and hash are in hex-representation.

The output listed in Tables 2, 3, 4 and 5 was generated by the phc_demo program, which is also provided as source code.

```
Teh quick brown fox jumps over the lazy d0g
$0000000000000000000000015:1$114495a461ae8b1500e532aa7173bb26e1a43a2e06bcee934a2d9cb714eed262afdb600fdf01677f8a300a1cd30eeb8d3f3c0bde4b11698c213e3a2a202cd75e
$0123456789abcdeffedcba9876543210$15:1$e0b0218195ff7bec008baf8784b361594ff81a33c7e0ce10e4fd976c6da53235e4fd9a75546d84657263ead1478164b4d6b6389a1dc6371705bbd342f281d51e
Input
Output
Output
Output
Parameters
                 $aaaaaaaaaaaaaaaaaaaaaaaaa$15:1$e49c68ebdc8fdc65e31a68bb7d3b241ed8c52dfa83ef539df76a3c0a877aca090568cab2273c7bf6dc6822f8ec852b96f0ea28e3b3339985de0660681a606e5b
                 \text{Value and the properties of the properties
Input
Output
Output
Output
                 $0123456789abcdeffedcba9876543210$15:1$1467b73b0a48be9beed57e3b7388ad9d17fca2b700e86dda38522bac2b59ac26d03ffadc9074528eefcbeb31115a5f17222d2639eaf68da22f1128b2d8e2641d
                 $aaaaaaaaaaaaaaaaaaaaaaaaa$15:1$708709296f3b2a870828b19cfcc3be3d2554dalec2lda46e35a8cde97ff418a320e9a8606a1539439623fd524c0497666f32d0c8e6d020807289f76094d0531df
Parameters
                 t_{cost} = 15, m_{cost} = 1
                 The quick brown fox jumps over the lazy dog
Input
Output
Output
                 $0000000000000000000000000000$15:1$8729969794ca741ff9c88ef6e394a61aedb20dbe407d464dda9a68386b3fa653eb03e5818314e636719c6f65bc88981057e082a48c2190f776b91eaae0cf51e4
                 Output
Parameters
                 t_cost = 15, m_cost = 1
                 Input
Output
Output
Output
Paramete
                 Input
Output
Output
Output
                 $0123456789abcdeffedcba9876543210$16:1$47fec4d92c8ca64a1fca7fa7d5fd446f213f33905c8669c27c6cbfba2681e55db5faf201c5ee62b95f9696ba6241a7dd8b7e8591594fc5b35a507fe9c673d9ca
                 $aaaaaaaaaaaaaaaaaaaaaaaaaaaaa16:1$3773c5483a204d5a0644ccf3bfc8890d53bfa1c877bbe29553e57a3acea179fa18fd4acec582a469dbe6af340bf86993d5bdb60e07df7d9a172ee4be4764f135
Parameters
                           = 16, m_cost = 1
Input
                 The quick brown fox jumps over the lazy dog. $00000000000000000000000016:1$c784ce408ae67fd285a0ce35705dc6663b699fcbb03ed7c1eca32b702dadeaf2b16b5d4b4c9e53e336fd86801e0b0174e409a123b402f0405233dc9cb58b0938
Output
Output
                 Output
Parameters
                 t cost = 16. m cost = 1
                 The quick brown fox jumps over the lazy dog $00000000000000000000000016:188ff883aa95752dedbf62ac6fef24dc7a4028d55478ffe739ec05b241c94948a783c01ab52f766039694e3a358a4c91d40b1f7f64f1fe4a93cf55431b5725ad5e $0123456789abcdeffedcba9876543210$16:18159a8b2395dc575feebf8783044483288901c95ba30fd925cd5bb1a5ea7604bb7a228c24e500d70ce59a0fc076c1cfdabf2a8179189f848746da50fdbd1b25d $aaaaaaaaaaaaaaaaaaaaaaaaaa516:18458db1cce4e011eecd281880955acf6797586feed20ef3589a030957a08be6b05d73c434eee7a2a840901ea035299a663b0db6a8f2db585dc3687dad5878ef2f
Input
Output
Output
Output
Parameters
                 t_cost = 16, m_cost = 1
                 Input
Output
                 \$0123456789abcdeffedcba9876543210\$16:1\$17585fcebc3bf5e9904a08fd771d0db19bfe74a2fc386afbaf9d1e70ffcbc116db45de5eedd907b277eeaffb9adeb86999b5babed36053bec2c01ed276bd4688
Output
                 %aaaaaaaaaaaaaaaaaaaaaaaa16:1$8773e25937d828eb2b30ca4031fb8cada72131b8861f1ae05b15d59988e5768ed1ab73540dfb2cf3c221fa577c8548e0ffd61a05da81fe23e456a343bc81c908
t_cost = 17, m_cost = 1
Teh quick brown fox jumps over the lazy d0g
Output
Parameters
Input
Output
Output
                 $0000000000000000000000000000017:1$3994d3a1f1b8e8c9f892e15cde440f87879193b66f6fb095d36e61718264e9eb5ca7e3de537aec41e89f516d63f735264304ec22fa8830134bf5f7a77303d034
                 \$0123456789abcdeffedcba9876543210917:1\$17335781ffa365480e8f0ebca2de0ef36330d731a23190614b38b414b57dccc93c661570dc1de44d82eda60863e63acc04d6aac84defd650d30b4765c238b884$ $aaaaaaaaaaaaaaaaaaaaaaaaaaa*\$17:1\$f9e6722302999e22eae3599ce48e613f857229c2d5b052323ed84c2c0f9ec138c25b4de48c4cbbae42e4df7ca551af9516bb5819b84f7e00dbf5972fdd1f2fdd
Output
Parameters
                 t_{cost} = 17, m_{cost} = 1
Input
Output
                 The quick brown for jumps over the lazy dog.

$0000000000000000000000000000017:1$394441cc7f1222521e81cf5aa2d5be04b9ab195a3747a497402e4c3d6a86fa353ff582d31e0aa22e248970f3a40637e9d3cb818c4f58fc8d4bd8a91fd23c3158

$0123456789abcdeffedcba9876543210$17:1$0620fa8bf893d08d98ab07324e6d8d17c4009d5ed6c093c3f4a1bb4b9758be52f2dc2676fe638def2679d8a3ef439ba5e3342f20e49c9eb26234e06b7a91caee
Output
Output
Parameters
                 $aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa*17:1$e8540d6326e6fa62293c0af90b86937e29b18c4edc912deef557f26b68b676baa5cacd1e1cbbbcd6a19dadb292f1c695ccdec835ab4b8082efcec49eee89977a
                 t cost = 17. m cost = 1
                 The quick brown fox jumps over the lazy dog $0000000000000000000000017:1$98c5c1ef349c231841aa88a8848aae1fe302546ac8a960c169597412e2c48f42c9bf86c21f5c924e0d0e88d7b786f4ba175c3ebffd04d21e0e34596406078917
Output
Output
                 $0123456789abcdeffedcba9876543210$17:1$fea87c4a381b520e66876246202a3b8c7bfb79c79b1c15b686f76c78c3d6e2bf6a81600f78c5cae343f8d49d35b77d8d26048d4ffe8ed574f8e5ea5495f7f943
                 . $aaaaaaaaaaaaaaaaaaaaaaaaaaaaa$17:1$577962eb40181e1daac38acba37568c9ec02da6d27ebdf0b2616b859717f61238fed5fb91c99c60df7adeb549c545105f7a26c23c1d7a3af8e591e0830415bd3
                 The Quick brown fox jumps over the lazy dog
Input
Output
Output
                 \$000000000000000000000000000000017:1\$778be27376cc46e81676b014322c6f69ca558cf729368b34a9862d012ec8fbd480a0b3cfcaa9adf63da515e8c116ebdae87ca5831b99bca62298b38f52106452
                 $0123456789abcdeffedcba9876543210$17:1$c1b51162b8f5a8e815938c5b15088e94eaf9dc14108957f1068397e5a809409228548841edb769228984893602af548e1cbd807335ca93ee23ad78acf77bafcb
Output
                 $aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$17:1$22842c532b6847e046619f54e851b989bef9a768fdcf128234cca327a38714623afcd4be226bdbf718de48fd6ffb329a0bb13cbcb1010bcb12c89c99a327b80c
```

Parameters: t_cost = 15, m_cost = 1

Table 2. Sample testvectors for multiple $m_cost = 1$

```
Parameters: t_cost = 10, m_cost = 5
                 Teh quick brown fox jumps over the lazy d0g
$000000000000000000000010:5$1b4199260f4b15730799253350c4a2b8e4d216ff2d820635c7a4076576cfbf458608e77060c8f399efb83fcf0e2a72fe742f6d337cebd51cc485ac5880c0bbbb
$0123456789abcdeffedcba9876543210$10:5$6ced729c917f3be9e66a081626878b947d4cb7569c212e75068b9ef7be761ac22367aafd4a1ffa60e2653a5865b40a8098bef1908872ab29329d0489ed2e87d2
Input
Output
Output
Output
Parameters
                 \$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa\$10:5\$d38f1d123371db95bb7c336d89ebc0b79c05e170ab46ef625cc4e5f054f6d70315c87f3bf8a0acfa39c35174baf93d4913f55eb5f384811b822cee95e2f2bed8
                  t_cost = 10, m_cost = 5
                 Input
Output
Output
Output
                 $0123456789abcdeffedcba9876543210$10:5$2f29a614354f6cd85843ae67a7a27d5c05b2e8f26b1d6237ef3c6e06d594b04594894f15d1b5831a7ed783f19b445249f3bb0389e062a4a24ae5b7fff764ccc1
                 $aaaaaaaaaaaaaaaaaaaaaaaaaaaa$10:5$5dcbc80755ca543f6b706f0c93f83fcd6de25ccb6e63d719f8546a1b06da4f13dc4c628df19c2badfba298d7e0522e63bc1a7b4bfee8ef9c7dd44234e37ed32
Parameters
                 t_cost = 10, m_cost = 5
The quick brown fox jumps over the lazy dog
Input
                 Output
Output
Output
Parameters
                 t_cost = 10, m_cost = 5
Input
                 Output
                 Output
Output
                 Input
Output
Output
Output
                 $0123456789abcdeffedcba9876543210$11;5$21f58def934627e1e19c573bb2b4f2b4b6c8df878e12f1907ea935ee194033741a416d63be405168729799d4f753379fbd4d12f1673587e61cd3db856e903e19
                 $aaaaaaaaaaaaaaaaaaaaaaaaaa11:5$4eadfc3b00f3bb82c1226476cb71b530b891c47879c87b246df6aa3c72495f6d685f1d7c9ff2379a89e1e982c53784079083eda2ca4555be5bcc5a22749f40d3
Parameters
                           = 11, m_cost = 5
Input
                 The quick brown fox jumps over the lazy dog.
Output
Output
                 \$00000000000000000000000000000011:5\$1b2a70c689f793de7a3a32ffe051733a6bc98c82fb5fb7eada538163ae3445fc1d7c012bf6176458add4f597be94c45bd3c629db8056043a7b99ebf0c25271e0
                 $0123456789abcdeffedcba9876543210$11:5$2be5b1dc6c292d1990cb6f646228c4ae756c1d8117aca6abdf78fd9244a00e465d29c4455a896d62ebe9f09beb0f0891a07a46d98f2ac93015dcf376cb7e7dd3
$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa11:5$0c18b345eec78fa9621f1a17bdf3641bbd516ddfd670449d0c6f53a91b9b03f2cfe25393e56b69229f3b8aed19805656e235ea7d5c5b6dd899eefbd72035d263
Output
Parameters
                 t cost = 11. m cost = 5
                 The quick brown fox jumps over the lazy dog $00000000000000000000000011:58ah648773078abf2cd7d2e8a69e3f850268e06e9faff9df02cbf7ef299a4a0c0d4f328b20041b1b1e25282eb6e88b97f78e9f345a6117e9e94af35eef3d67371d $0102456789abcdeffedcba9876543210$11:5816ff685274365c3c97342c52d554dd437e96e24d1bdc9ba699fc1a8e1909b0f4cab631ece6668fa97431f3c5a00bf3a7b93bc329a6cb69b88c892692bdb2fd0cc$aaaaaaaaaaaaaaaaaaaaaa1:1:58768ed5374c5552547015fda7e9e3c9283dd9cff724b7f7d6e4269ade3d665a7f65236fcde21bd8169d592f7b68a5e88e7785b4ef7c1a9439a4f23e05ad7b46a
Input
Output
Output
Output
Parameters
                 t_cost = 11, m_cost = 5
Input
                 1. m_cost = 0. m_c
Output
Output
                 %aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa1:1:5%a6e2ebe2f5fb59d6e761e56cea942333705f6149acdcfc570a904165dd6a36956b6ef553d4870701053d91ee6876f060b93300025a5ece23a75eeee9c9702c1d t_cost = 12, m_cost = 5
Teh quick brown fox jumps over the lazy d0g
Output
Parameters
Input
Output
Output
                 $00000000000000000000000000000$12:5$b2a8f5c30fa48ff210e219e431f61f6f6ae0f98aafe3e6ed4ed53f1081b7cd94622421ae366173737a0245b06cd115a86e29b398a0c0b0ffcfc7ad11be981858
                 $0123456789abcdeffedcba98765432103121:58aa3abc18c5e7e361a42e5302264e45659e22bd29e5550e74a1e1de76c7c3863877720c24ef072e74dd22ddc77e2fba3f799dc8602b86e87fb552f1a892aca$
$aaaaaaaaaaaaaaaaaaaaaaaaaa12:5$bc22ee2bdaeacf50908338fbcb1eaefab7b7aa4083edc225d9857fe84661a3966c3c03659ab937a7a090fb313102f2395acf013872e71686b7b6ceddf048ef69
Output
Parameters
                 t_{cost} = 12, m_{cost} = 5
Input
Output
                 The quick brown for jumps over the lazy dog.

$000000000000000000000000000012:5$44b2963a0d9fb862c3bfa20c4bc34e21439f7d20576d16d0047a5b6bdae681a9765434f9cf5754f807be6663446d7550ea59abd8660b0c6bef03b24ce34c9863

$0123456789abcdeffedcba9876543210$12:5$4602b526abb0f7541e49b928838bc40551e7bb5890c42d6c19ac9c5f6c250e20f741e9fd52c37c9788a6a702565ba448af22f90e389ffb430958db361fe1b9b5
Output
Output
Parameters
                 t cost = 12. m cost = 5
Input
                 The quick brown fox jumps over the lazy dog $0000000000000000000000012:5$e464bb24817858b56e4b81fd2dc2226b7abb2b024671701e3b2f6ba7c3fd9f2d0bb58d7b2e98bec29467d590e5d1f371050c700711a452bb19c195a709d9ed0c
Output
Output
                 $0123456789abcdeffedcba9876543210$12:5$7198bfc51d32abf7a8893b3ad7229d98a34efe656c62788493a688784a13fd02722ef19f3873da1cc7753084ec78f65c9a45e259bc45856c6828cedd065eddfb
Output
Paramet
                 . $aaaaaaaaaaaaaaaaaaaaaaaaaaaa$12:5$49589723a3745ce463d312c885efaf1503b1e84c80740804867c8a4891e9503d9faa2e28694a8b95eb2c7bedbaf4cef06ca1e873b861795e11031034f2811dca
                 The Quick brown fox jumps over the lazy dog
Input
Output
Output
                 $0000000000000000000000000012:5$9a0c560f421f6a940b7201fdbe0404d980b530cd0a6266a18d971d0f7b6e3a93ffd1499f56a879329b8650010cf604003bbefd3a012c319edf3bc8a93ad6b676
                 $0123456789abcdeffedcba9876543210$12:5$7846747960b90eb740b133c6f81b080300e4b492a2f1eafa0986659186c248b357fc46b61aaa10c2b3dd8b4224466b7f13eb126243a7d23ebccbbc53603a511c
$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$12:5$f0e6c5b774edc3708fb3f695fa9ef12b90bda9a40779eccda62d6b0a0cac06b948fe10f448f67c636595b2390f3cb4129f89967d3402116263e4ce78e56bf672
Output
```

Table 3. Sample test vectors for multiple $m_cost = 5$

```
Input
Output
Output
Output
Parameters
        Input
Output
Output
Output
        $0123456789abcdeffedcba9876543210$5:10$4eddb86e96a8c6d3f4eb1902b46a0c6c4d1c6f26d9fb2c4ea17deda85867f7b4a47b50d30d6adf68492de5fdf7028c7fa7508ac1dc46a12d89669ec5d82b98a9
        $aaaaaaaaaaaaaaaaaaaaaaaaa$a$5:10$d9f0b3ecc105c6d6478b5980d5e5a51ae46f73f34fc0616a3bc7d4e2cd9a5d3e15af4592ef80d4061bccf815b891d394402eb6c44fbef0859fc69997c7ae65
Parameters
             = 5, m_cost = 10
        The quick brown fox jumps over the lazy dog
Input
Output
Output
        $0123456789abcdeffedcba9876543210$5:10$6a2df69431fc1aca0723e7fa6c7729f8d97d0b9eaabc1968a917b62661c2477001df3f5e4f4eb80bf4cd4bf912033360ac606f111c3e8cefe9e0330dd6e091892$
$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa5:10$fb7e62a0ffa01646aa3b2bb251e2628c0bb4ec0f9e219dad27342b104ed918f91f8ac6edb8ee5d81585cfd8f705b232021519719b49bd000d5144c4544ac83b8
Output
Parameters
        t_cost = 5, m_cost = 10
        Input
Output
Output
Output
        Input
Output
Output
Output
        \$0123456789abcdeffedcba9876543210\$6: 10\$67af23812c8c59c52121821876b3a0323fbb7b7a693d84126b7bf1fb45e87d50f982e2167fdfaefbd56549bc128af254e4da42180cec4f87f31ef4b5d25b74d0
        $aaaaaaaaaaaaaaaaaaaaaaaaaaaa$6:10$1591c133c6c7f33f62ae17da8c412afbd8f6c2676e24ad7ae80778ddc88429c5bdf55b67c7a2b56e6994510ebee41392477c4f162a708825e968ef99cbb14777
Parameters
Input
        The quick brown fox jumps over the lazy dog.
Output
Output
        $0123456789abcdeffedcba9876543210$6:10$6666ef422a9d677018929454ca333239cc477d9696e055d45ff3364d6aa91ba6330ff2b3ee1ac859b864d08186417f82ade1e98c756b3acbf431687c649e83819$aaaaaaaaaaaaaaaaaaaaaaaaaaaaa6a6:10$cf9e3a8e1e597ee1e122a32c81953274e3a95b2c0f20fdbec59ed6ee550ad718b0c432d2f86ee2b8df3caebb191e32a68dfcd990ae97e2711702ab652ef4d8e7
Output
Parameters
        t cost = 6. m cost = 10
        Input
Output
Output
Output
        Parameters
Input
Output
Output
        \$0123456789abcdeffedcba9876543210\$6:10\$1a7ff28e76ad8f4d63888991e4745fcef9632b7ef5cbf30b7d9edac0d6ea14031a0f40ffac6478aead8151fa908b9e14d32717a0cc215239e7ec792cf4a9ad
Output
Parameters
        $aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa6:10$f7bc33efc3f47498c33fca3e2f1ae838a6810de3caa07c553631841e408ad720c83550dca23eafc236d99d488080c4b9a663d71204ac015662b79074cc1982ad t_cost = 7, m_cost = 10
        Teh quick brown fox jumps over the lazy d0g
Input
Output
Output
        \$0123456789abcdeffedcba9876543210\$7:10\$16181b1c33199b5fc4edeff7a732dbfe031d9b737982154c9e83de32a2043f1981983c1f5fe4decff9a6faf223f145a841e43f34aa94ba51bbde3c6e87ad08c6. \$aaaaaaaaaaaaaaaaaaaaaaaaa*7:10\$a05bedf903233a09e892d9783479cffa28e74d7f091d4bdf17399f2032696b0896dda55f7fb58382543a4e9627e7a4f7828e39ed112d5635d984767178a533bc
Output
Parameters
        t_{cost} = 7, m_{cost} = 10
Input
Output
        Output
        \$aaaaaaaaaaaaaaaaaaaaaaaaaaaaa^*7:10\$458715aa62abc0767ba9f4efdcd4d5c4c9359f4caced2ce8f1bbed65ada5b240a66851578a2288c404db1c4ef6ed5707fcc8831945aaa51aaf4b16a80fc352e1
Output
Parameter:
        t cost = 7, m cost = 10
        Output
Output
        $0123456789abcdeffedcba9876543210$7:10$a8f4e8569f8231ea54897e207ba7aa9c69a524102b11b92c50127d93ea09ae5e18de5094ce097b6583db89c5efae7a3edd401668efb17a171d725bfc0378e3b7
Output
Paramet
        The Quick brown fox jumps over the lazy dog
Input
Output
Output
        $0123456789abcdeffedcba9876543210$7:10$760cea496193d665af62c46a1ff1eb5cb88625292dd1a2afaa58e0550cdccae596fbb235d62f1f108bc66fbee93217bcfffe8124f3cb2737b33fffa50c908f49
$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa*7:10$59fecba279f016f54235d60dc543c1ef3e18752c25a06b9a4607a97bb66898728c17e8fb987abee08a6ca16e2c5f5b49224c66dcd568f7787b9ea890d4ac023b
Output
```

Parameters: t_cost = 5, m_cost = 10

Table 4. Sample testvectors for multiple m_cost = 10

```
Parameters: t_cost = 1, m_cost = 14
           Teh quick brown fox jumps over the lazy d0g

$00000000000000000000001:14$4ceb63b51cd1b6bea96a7347cd9dd3b8d1e0f31d600214249ef17c872ee3a18847e59503c59fb6bca3add84b2ccc3f7625703e476e07d6690951f1c619cd9a7b

$0123456789abcdeffedcba9876543210$1:14$cf4b36b561374a45782ffaa8e7a20f52b77897317c6f6cd5a1555a5d1ecd78cce345def6c76494c857ea7f8a004018977b5238d04a3718a70f7d565e3de26504
Input
Output
Output
Output
Parameters
           \$aaaaaaaaaaaaaaaaaaaaaaaaaaaaaa*1:14\$0982c64f4730c29dce65c48155ff82534c4f0f90e215415f9d67a5f829cd00fca44ba84681d5f326256e7ffa43b05727b923c9b015d463f766910ccdc9044d54
           \text{The quick brown fox jumps over the lazy dog.}
$00000000000000000000000011:14\text{cost}=6229\text{boso}104\text{boso}18578\text{bbca}26625\text{pag85}ca5c6809a26f0a60994076ee18b649371f5073eed9f9b37864204fe95dba7c12ff6423c34cce968a2c5dd}
Input
Output
Output
Output
           $0123456789abcdeffedcba9876543210$1:14$76b0ebf36a1025cb37ae01a813a4a4b60b3a40b09e22fd862a3b606a80b04be65046b8bfcfb0b690634785ac0f99dbba032d8df0ca02a6e5ba73bd72b2f
           $aaaaaaaaaaaaaaaaaaaaaaa$a$1:14$cd78e1c22c4da74264f3f1c7b62a0772adedd873d40716be1cb50d5e37cfbabd3856d7dfdcae3bda8a28ce376788db1bc1e2c0a33a96bbc47fffb542d01a65bb1
Parameters
           t_cost = 1, m_cost = 14
           The quick brown fox jumps over the lazy dog
Input
           Output
Output
Output
Parameters
           t_cost = 1, m_cost = 14
           Input
Output
Output
Output
           Input
Output
Output
Output
           \$0123456789abcdeffedcba9876543210\$2:14\$e714c86f80c58ab537f400a368ba987416176a2394017fa842c15b5a438e76f77012f26b04da1f3b02e802f5d3d3cf81eb3590ac73047a3806812345137cff1d
           $aaaaaaaaaaaaaaaaaaaaaaaaaa2:14$40afb9ea973eb2c1ddcc03eb575d7b3a8e446ab3110706cbcfdc820d16debe7e17b1fbb2f03d8ef20eedde72c0bbdf2dd0b76d95ca37b17afe1e2c6244809bab
Parameters
Input
           The quick brown fox jumps over the lazy dog.
Output
Output
           \$0123456789abcdeffedcba9876543210\$2:14\$87e04c842d1091226a0090e2597c4b6:14303b245039ea1b021796fa2aa27307c642e77c412d107ba11fde25061a1d20f448222bea195174792d0638cbfbea66 \$aaaaaaaaaaaaaaaaaaaaaaaaaaaa2a2:14\$9407239a6f3522489c329a59c04706ed94a47cfdaa8cf9885ce625cd77a9d7b2368451e51c7d04fab51666c93054b4880a3c1c7a19836b9e61b121ab383b6451
Output
Parameters
           t cost = 2. m cost = 14
           The quick brown fox jumps over the lazy dog $000000000000000000000000010012:14$edf2229e6c3017d985c1f7c488ca2e4e339c772a26e72965b22d91e1969796f291942a24f1b2fa9f0b6dce0f6f35e643330fff53dace99d05acd1979287e738a $0123456789abcdeffedcba9876543210$2:14$bc758aabe50f40690eb1d6886d4a29cd9d10a681fb434718f148cafd25208ccd6244b5e1239d5fa2fca5858b22f33ac2f668d5c2ca07ff72c3d7c64d2a258b71 $aaaaaaaaaaaaaaaaaaaaaaa2:14$b07556d41802e053cd112c2bdab2fd10598e52ec58b616e51c52e42ceafa5616b1d15fed14d24e29390ac91e37feb63f04132a4393f5f14d74703283f33c81b91
Input
Output
Output
Output
           Parameters
Input
Output
Output
           \$0123456789abcdeffedcba9876543210\$2:14\$b57516b2c98ee6c32cf24b1359167f3b328790263c23260d9ed373e235f6ca3869b886b34f6ec8d7ff0fa18d416ec97bfc3b0ddc6e20e438b365e50bdde44909
           %aaaaaaaaaaaaaaaaaaaaaaaaaa2:2:14$69110d775f9189b18826535e4a46639dd93a245851bf8c006816a1f736e4a6e736da3f64e677e64a3fa4ac7b19b2c748872e101cf560f0d00ebd4d86e3a9e3c8
t_cost = 3, m_cost = 14
Teh quick brown fox jumps over the lazy d0g
Output
Parameters
Input
Output
Output
           $000000000000000000000000000033:14$16174c893e96cfb110af584caee53bd589ef1d99c87d1ca512a6646383e01c1cefcd176f4b261d4106300b861fcdc64a9b6a92d5e06d441c262e0d7dabb31a25
           \$0123456789abcdeffed:ba987654321083:1487794ddd898c5c29c7a347216079a07e90b39533464473821545118c0ccde91d8e0b88b0b7906e38379454193192831c332995010a41717a89ea035ede7ddec \$aaaaaaaaaaaaaaaaaaaaaaaaaaa33:14\$0bab08aaa6202226c6a7f9b7f429f48821742e7ad6c227fa329fc7a8c5ecaa5d13feee5e54cbb766d79c5be24f2bb64b25207b2cc096e1b6e3370f3c3725c428
Output
Parameters
           t_{cost} = 3, m_{cost} = 14
Input
Output
           Output
Output
Parameter:
           t cost = 3, m cost = 14
Input
           Output
Output
           $0123456789abcdeffedcba9876543210$3:14$ac133de45f75179daa1dffd6a6e56e751deeb36d17c2e736ef9141a6157baeb90ded8c904f18566bc4b307a7bfd67842307cd1eb624e0e05e6be25e079075c80
Output
Paramet
           The Quick brown fox jumps over the lazy dog
Input
Output
Output
           \$000000000000000000000000000000314\$fb956fc5f9016e18bdc98c7bc55bc2c7f953ce70ba9b99b5b5da93a6d087e93a83158707cfc02262928cd0f3a816f28d3311c7c7f4973284a199edc5c0c9d84e
           $0123456789abcdeffedcba9876543210$3:14$ec29f7508697679ae2034d240e7897a1cb6016e1bc35e40725265e01b087b55b3f7cca4715d86e176df3c220a36aba2d01c26033f7f700cd819fa19529032c6e$aaaaaaaaaaaaaaaaaaaaaaaaa3a3:14$b57d6567eb16b819dd9c97f2574fe82e0095cc055a7b58fddb560a3c9e6c0ab14b1251723ce886caefe947b119edd7c2316357296e347f1e366913e42756910a
Output
```

Table 5. Sample testvectors for multiple m_cost = 14